



**HAL**  
open science

## Modeling an IP network for audiovisual streaming to improve double failure recovery

Stéphen Pirlot, Eric Gnaendinger, René Kopp, Francis Lepage

► **To cite this version:**

Stéphen Pirlot, Eric Gnaendinger, René Kopp, Francis Lepage. Modeling an IP network for audiovisual streaming to improve double failure recovery. IEEE International Conference on Advanced networks and Telecommunications Systems ANTS 2015, Dec 2015, Kolkata, India. hal-01237339

**HAL Id: hal-01237339**

**<https://hal.science/hal-01237339>**

Submitted on 3 Dec 2015

**HAL** is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

# Modeling an IP network for audiovisual streaming to improve double failure recovery

Stéphen Pirlot<sup>1,2,3</sup>, Eric Gnaendinger<sup>1</sup>, René Kopp<sup>3</sup>,  
Francis Lepage<sup>1,2</sup>

<sup>1</sup>Université de Lorraine, CRAN, UMR 7039  
Campus Sciences, BP 70239, Vandœuvre-les-Nancy  
Cedex, 54506, France.

<sup>2</sup>CNRS, CRAN, UMR 7039, France

<sup>3</sup>TDF

TDF Metz

1 rue Marconi – 57070 Metz, France.

stephen.pirlot@univ-lorraine.fr

**Abstract**— Protocols engineering of IP/MPLS networks are constantly improving with new separated features and new resilience mechanisms. In the transportation of audiovisual signals domain we must compose with multicast protocols which are designed from other scientific developments. This audiovisual traffic due to its non-elastic nature presents a very huge sensitivity to network recovery after a failure and these effects can be amplified by end devices (encoding, decoding and MPEG IP encapsulation). In this way when we choose between engineering solutions the unique criterion of availability is not enough, we must complete by an impact analysis on the service made by the network resilience technics. In this paper, we propose a first approach to analyze the behavior of different protocols engineering to improve selection. We propose using Bayesian networks to compare performance on different criteria and we will illustrate with two engineering models. The results focus on a real improvement of availability by choosing the adapted engineering solution.

**Keywords** — Protocol Modeling, Network Survivability, Bayesian Network, Engineering Choice, Availability, Dependability, Multicast Transport.

## I. INTRODUCTION

A network service can be implemented using different ways considering topological and architectural views of the network, considering protocols accumulation and obviously considering quality of service requirements. Otherwise the client requirements are rising while we are observing a clear fall of the performance about our transmission links which constitute the network infrastructure and especially on the service availability point of view. Then we must strengthen the mesh of our network and improve our protocol engineering solutions to use this entire new infrastructure for all point of service [1].

Networks specialized in audiovisual broadcasting possess particularities which are not a prerequisite for an efficient telecommunication network. Indeed audiovisual streams are non-elastic and continuous so a short link failure induces a display perturbation for the final user. In this context customers are people who receive the TV at home and clients are TV stations which pay for a broadcasting service.

The requirement of all networks is to be the more stable and the more resilient. In this context where we are talking about network survivability [2-4] the problems concerning failure resistance or cyber-attack are primordial. In addition we want that all recovery decisions of the network are completely automatic to react as fast as possible to a failure. Most of the

time engineers have different choices of implementation for the solution.

Network services are implemented for each client and it is possible to adapt all the protocols for each client demand. Even if survivability of concurrent services in overlay networks has been studied [5], we are concerned about one specific service and its choices of implementation.

Because of the innovation of technologies and all limitations, engineers are encourage to evolve engineering existing solutions to new innovative solutions which are bringing better performance. In the multicast transportation domain the protocols unified under the “multicast VPN” term have proved their fullness and are now available with different variants. The choice of evolution and the comparison between two engineering solutions is most of the time conditioned by the experience of systems architects and with all the technical data given by parts manufacturers or by realizing prototype. The aim of this paper is to propose a decision support for the development of an engineering solution specialized in the multicast stream transportation guarantying all the above criteria. The first part will expose the modeling problematic and the target of this study. The second will explain the two engineering solutions, then the last part will present on simple cases how to apply this modeling solution and will show us on a real case the contribution of each solution to conclude this study.

## II. SERVICE MODELING

Availability study for a client’s service is based on the physical infrastructure analysis but it is especially based on protocol usage. Availability of the infrastructure bound to the service is obtained from the availability of each component of the network, for our case routers and links in optical fiber or microwave and from the constituted topology of these elements. When it is possible the real availability of these elements will be used.

There are many approaches to model the availability of an IP network depending on the physical infrastructure like fault-tree-analysis, Markov chains or Bayesian networks [6-9]. Measured availability will be the same for each used model. In this study we will concentrate on Bayesian networks.

Very used in dependability field, Bayesian networks are probabilistic models utilized to focus on some precise characteristics of a system or a sub-system. They represent all the functional probabilities of the different elements and their

interactions with the general state of the system [10]. They are presented in the form of directed acyclic graphs. This kind of model is used by system engineers to model hazards of complex systems as automobile factories, nuclear power plants or fighter jets. Bayesian networks allow predicting the global system behaviors to diagnose the reason of a noticed phenomenon in the system but also to control the system behaviors.

This way is the common way to complete a Bayesian network to understand the system's behavior. At this point we remain at the level of a fault tree analysis. When we are choosing for engineering solutions it is quite simple to know which case will have the best availability by analyzing all the possible failures that each engineering solution can deal with. But modeling the solution must give additional functional parameters of the network to qualify some eventual loss of performance.

### III. STUDIED PROTOCOLS

With this study we are looking for some decision about which engineering solution provides the best quality of service basing on theoretical data. In a network dimensioned for audiovisual transportation technical constrains are strong on certain criteria. In addition we are talking about broadcasting solution through an entire network to reach services' points separated more than 1000 km using hundreds of routers. For everything working we use multicast streams because the client wants a point to multipoint service. Today we have at disposition different kind of transportation of multicast streams over an IP/MPLS network and each solution has his advantages. Historically it exist an implemented solution over the network for broadcasting national programs, but it is planned to change for a more recent engineering solution.

The first solution is based on a ring protection (as RSTP solutions [11]) and is called imbricated rings solution. The second solution uses multicast broadcasting solutions and is called multicast tree solution.

#### A. Imbricated rings solution

This solution uses imbricated rings architecture (main ring S1-S2-A-B and sub-ring A-D-B) and on each ring an own ring protection is implemented (Figure 1). Routers behave as switches in this case. The main ring is managed as a simple ring by RSTP. The sub-ring is treated by RSTP as a ring because the protocol consider that the link between A and B is always operational and so the RSTP blocking port can only exists over the sub-ring A-D-B. The stream is generated by the two sources located on the routers S2 and S1, and it is broadcasted over the entire network using the topology made by RSTP (the cross on the scheme show us the RSTP blocking ports).

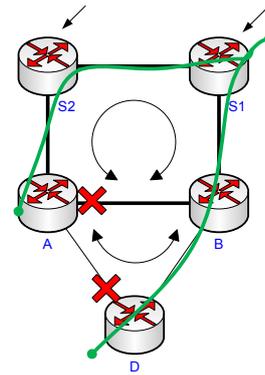


Fig 1. Broadcasting sample in the imbricated rings architecture

In order to protect the network from a global congestion, engineers conceive this solution with special filters on the routers which join different rings (here A & B). These filters are blocking broadcasted data circulating on the sub-ring to supply the main ring limiting broadcast storm phenomena to the sub-ring only and not over all the network.

This engineering solution is quite simple and is well dimensioned for broadcasting a stream over the global network. It reacts rapidly to simple failures but possesses some blocking cases with double failures. The Figure 2 explains two blocking situations.

In the first case there is no failure in the sub-ring so the sub-ring's protection is not activated and the router D can't get the stream. In the second case, there is no failure on the west ring so the broadcast is blocked on the router F.

In addition this architecture in imbricated rings allows only the ring usage and so there are only two injections possible for a sub-ring.

This solution is exposed to a recrudescence of simple failures and especially on double or triple failures. To find a solution, we must reinforce the architecture with more interconnection. The imbricated rings solution cannot completely take advantage of this network densification so we opt for another solution, the multicast tree solution, which solve these failure cases and clearly improve the system availability.

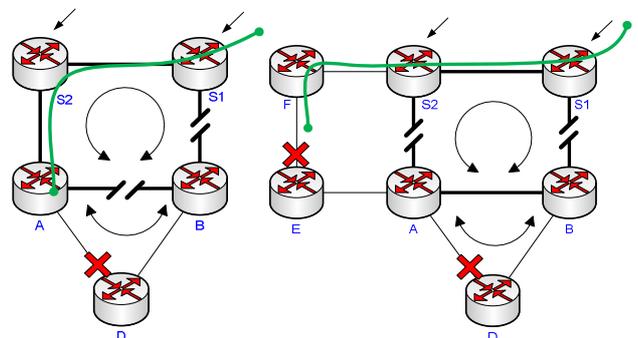


Fig 2. Two blocking situations for imbricated rings solution

### B. Multicast tree solution

A protocol engineering solution is always studied in relation to the physical infrastructure. Then this solution lean on two functional layers: the VPRN (Virtual Private Routed Network) [12] layer which emulate virtual routers and the VPLS (Virtual Private Lan Service) [13] layer which emulate virtual switches. The dichotomy is historical and technical because only the most powerful routers (nearest routers of the sources) can use VPRN features; the others will use the VPLS layer (this represents about half of the network components for our network). Figure 3 presents the dichotomy and the used technics to transport the multicast stream.

The multicast stream is broadcasted in the entire VPLS from a single and unique virtual router. For the VPRN side the PIM protocol [14] builds a multicast tree based on the multicast demands to distribute the stream only to the routers with a connected client or to the routers connected to VPLS which require the stream (using IGMP protocol [15]).

Furthermore, VPLS can supply the stream to the VPRN layer to avoid the isolation of a part of the network because of a misplaced double failure. So as a physical solution exists to distribute the stream to the destinations, the system will work.

Then both of the cases presented in the paragraph A are solved by this new engineering solution: it represents an incontestable gain of availability for the system.

Anyway from the functional point of view all the routers from a same VPLS (so all the members of a sub-ring) will be supplied by the same virtual router. This represents a true weakness in comparison of the historical solution because the blocking port is always positioned at the half of the ring so half of the routers are supplied from a side of the ring, and the other half are supplied by the other side of the ring (Figure 4).

The system can be weakened because many routers in the VPLS side are bound to a unique router which forwards the stream (in this situation the router B). Admittedly all the VPLS members will topple over the second router (here the router A) but all of the points of the service in this VPLS will be impacted by this switchover.

We are focusing on the compromise to establish between the gain of availability and the regression in term of performance linked to switchover mechanism.

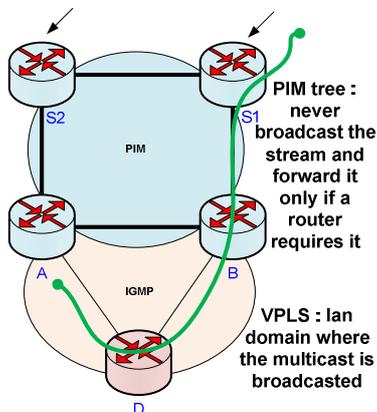


Fig 3. Recap of the used features in the multicast tree solution

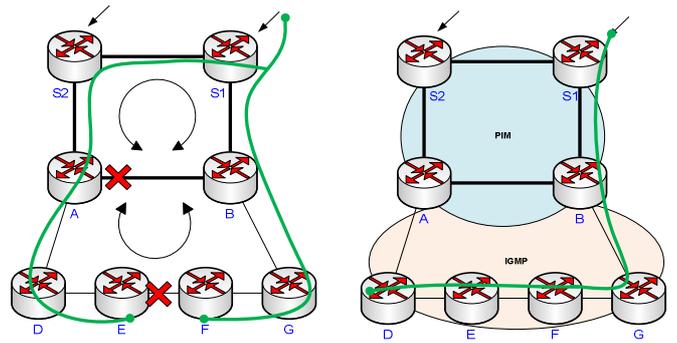


Fig 4. Broadcasting divergences in the sub-ring

## IV. MODELING SOLUTIONS

In this study we will focus on the system's behavior and we will look for modeling global availability of the system with Bayesian networks. It is possible to create such network by using the physical characteristics of the components. Links represent dependencies between the different variables (nodes<sup>1</sup>). The nodes of the network correspond to the random variable used in the calculations. It exist two kinds of nodes:

- *Parent node*: which contain a probabilistic distribution (the value of each component's availability, theoretical or real, is placed here)
- *Child node*: which characterize random variables in the form of conditional probability table (the target node, will be for us the node which modeling the system's state, is this kind of node).

Elaboration of the Bayesian network is realized in order to faithfully reproduce the behavior of the two engineering solutions without taking into account convergence delay that can last some tens seconds. We will especially focus of the gain of availability of the second solution versus the loss of system's stability.

To model the fact that a system is more sensitive about rerouting considerations (or switchovers), we use multi-states characteristics of Bayesian networks. As a matter of fact, each node of the Bayesian network (or random variable) possesses different values, and the choice is not limited to Boolean values as "Working" or "Not working". According to different failures cases the system can still function and these functional conditions could be used.

We use three states to model switchover behaviors and failures cases:

- *Nominal State*: When there is no failure on the principal way the stream uses this node. This is the normal working condition.
- *Rescue State*: When a failure occurs on the principal way, the system being structured to resist at every single failure, it is still working but in a degraded state, implying rerouting. To consider this situation this functional state has been created. The system keep fulfill his mission in this state.
- *Failure State*: Finally when the system is subjected to multiple failures that avoid the system's mission, the node is in this state.

<sup>1</sup> In this paper, the node terminology always refers to a Bayesian network node and not at a telecom network node.

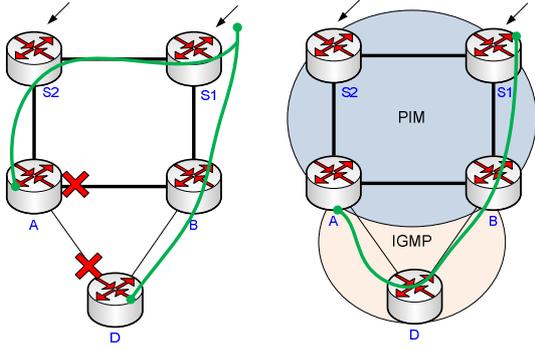


Fig 5. A first case to see the feasibility of modeling

In the following examples we will show that this state utilization will permit to quantify the loss of availability on the nominal way versus the gain of global availability. In other words is the availability of the system enough increasing to balance the rise of rerouting probability?

#### A. Feasibility demonstrating

Here we are going to show that this modeling way is feasible and corresponds to an expected behavior with obvious cases. In the following situations, the two sources S1 & S2 will have a stream at disposal and the studied router will be the router D located on the bottom on the scheme. The solution in imbricated rings will be placed on left and the multicast tree solution on the right.

The source streams which come at S1 & S2 will always be available and we will use estimated availability with the following values:

- Router availability = 0.99999
- Link availability = 0.9995

On Figure 5 we note that in the second solution, the VPRN side is not usable on the router D. The broadcast will be done by the router B which will broadcast on the entire VPLS domain (here only one router).

The stream follows the same way on nominal situation in both solutions and the new engineering solution will not degrade any performance. The only difference between these solutions is about failure detection. Indeed if the router B is isolated in the first solution, there will not have any rerouting and the system will be in failure state (cf. Fig 2).

When modeling with Bayesian networks the multicast tree solution as presented in Figure 6, we can find different elements as parent nodes (in blue) and some random variables created to easily aggregate treatment possibilities of the stream represented as child nodes (in yellow).

TABLE I STATE PROBABILITIES OF THE FIRST CASE

Solution	Rings	Tree	$\Delta$
Normal	<b>0.99897028</b>	<b>0.99897028</b>	0
Rescue	<b>0.001018685</b>	0.001019199	0.000000514
Failure	0.000011035	<b>0.000010521</b>	0.000000514

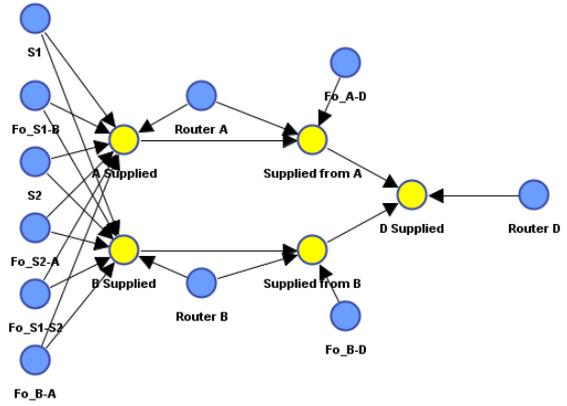


Fig 6. Bayesian model of the multicast tree solution in the first case

To simplify the analysis we will check if each part of the network is able to get the stream. For example the node “A Supplied” has for truth table:

- Nominal state: Never because the router A only broadcast the stream if B is not able to do it.
- Rescue state: If router S2, router A and also the link between S2 and A work, the node will be in this state.
- Failure state: Every other case.

With this principle we can complete the model and do the same for the imbricated rings solution by considering the blocking case of the isolated router B. It is quite simple to compare the performance by using the inference algorithm included in Bayesian modeling software. We get the working probabilities of the service on the router D reassembled in the Table I.

This model can directly show the availability gain by comparing probabilities to be in a failure state between the two solutions. It is a theoretical case with estimated values and that is why we are observing a light gain (approx. 16s/year).

Now considering that this model is able to show the gain of availability, we are looking to highlight what we are losing with an unfavorable case.

#### B. Study on an unfavourable case

In this part we are looking for a definitely different behavior between the two engineering solutions as presented in the Figure 7.

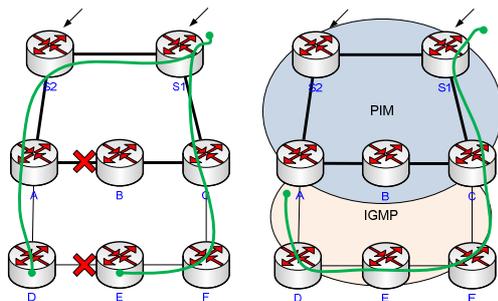


Fig 7. A second case unfavorable for the new solution

TABLE II STATE PROBABILITIES IN THE SECOND CASE

Solution	Rings	Tree	$\Delta$
Normal	<b>0.99846081</b>	0.9979516	0.00050921
Rescue	<b>0.001527375</b>	0.002037358	0.000509983
Failure	0.000011815	<b>0.000011042</b>	0.000000773

In this topology, with the imbricated ring solution, we can easily understand how the stream is delivered to the router D with the ring protection. For the multicast tree solution it is the router C which distributes the stream over the VPLS and then all the VPLS side routers will get the stream from the router C. Because of this behavior we are in a negative situation by the fact that the router D is able to get the stream in the nominal way by passing the routers chain from the router C. Anyway we still have a gain of availability because the new solution still can correct the failure case of isolated router A (which forwards the stream to the router D in the imbricated rings solution).

The results of this model with Bayesian networks are resumed in the Table II.

Here we have an interesting result: in this example the availability is improved but we observe a loss of probability to be in the nominal state. We win 24s/year of availability with the new solution on the 373s/year (+6.4%) that we got in the actual solution. Then we increase the risk to be in a rerouting situation of 4.5h/year on the 13.4h/year (+33.6%) of the ring solution. This case is unfavorable because by extending rerouting risks the performance decreases for the audiovisual network.

The next part will permit to quantify the performance of the new solution with an infrastructure based on an inspired real case.

C. Modeling a real situation

This part will show us how the new engineering solution will be able to be deployed over a real network and we will see how it contributes to solve double failures that are actually blocking in the existing solution. The architecture is shown in Figure 8.

With this study based on a real architecture, first of all we observe a gain of 164s/year which correspond of a 28.4% gain compared to the 557s of annual failure. On the rerouting side we note an improved probability to be in a rescue state of 24min/year against the 27.7h/year which correspond of an increase of 1.4% of rerouting. This difference with the precedent case comes from the fact that the main ring can be supplied by the sub-ring in the multicast tree solution.

TABLE III STATE PROBABILITIES FOR A REAL CASE

Solution	Rings	Tree	$\Delta$
Normal	<b>0.996824298</b>	0.996784525	0.000039773
Rescue	<b>0.003157427</b>	0.003202403	0.000044976
Failure	0.000018276	<b>0.000013072</b>	0.000005204

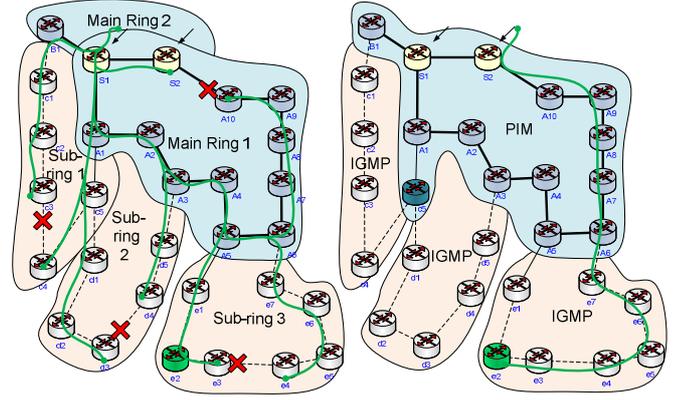


Fig 8. Study on a real case

This illustrates that in this real case the availability improvement is more important than the increase of the rerouting probability. In this way the new engineering solution works as expected.

This study quantifies the gain and the inconvenient of the new engineering solution. It can be extended to the entire network and already reassure the choice to develop this new solution.

D. Scaling the model

The study over only one service's point is not representative of constraints from point to multi-point topologies which are specific to broadcasting networks. However this scaling requires certain strictness in the modeling in order to never forget any specific case and avoid cycles (it is a characteristic of Bayesian Networks). Then some of simple and recurrent topologies of network can be modeled with a single model so we can get a translation from a network topology to a Bayesian Network. Some examples are presented in Fig 9.

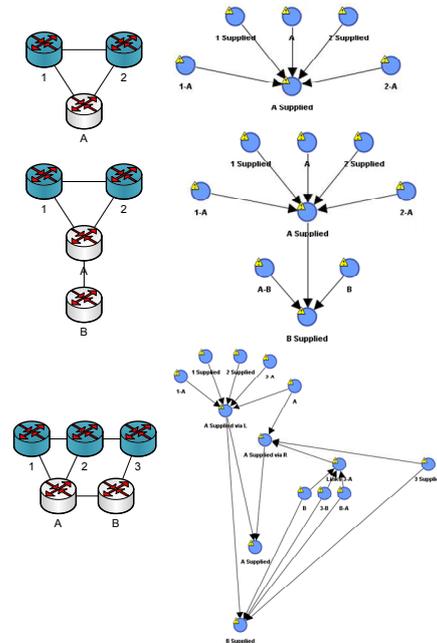


Fig 9. Some translated topologies in Bayesian Networks

## V. CONCLUSION

Multi-states characteristics of Bayesian networks provide us a deeper analysis than a simple fault tree analysis study. This paper presents its practical application on a network in a telecommunication context. For implementing a new protocol engineering solution we could quantify at the same time availability benefits aspects but also the compensation induced by switchovers situations. The results are in favor of the new engineering solution by introducing new routers with improved features, placed on an optimal way to reduce undesirable effects.

The future of this study will be generalized in that way to generally optimize the network architecture. This modeling solution highlights benefits and losses of each solution by using real data. We could imagine apply this technique to estimate the gain on other modalities, and precise decisions. With this study the instant switchover phenomenon was not taken into account even it is a problematical topic on audiovisual networks; the following of this work will be able to be concerned by this.

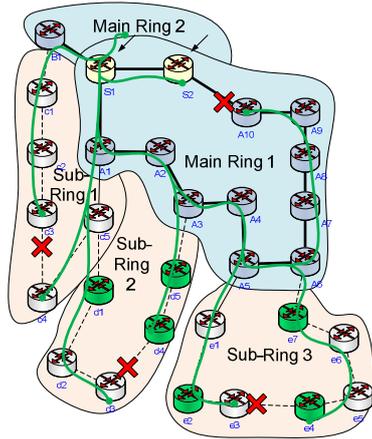


Fig 10. Study on six service's points

To model every point in the service we can proceed by the same way in the precedent case by simply adding a new node called "State of the Service" which will be the aggregate of all the points of service.

This solution allows modeling more complete and more representative networks of the real behavior of the service. We can apply this to the last real case as presented in C by using all other service's points as presented in Fig 10.

All the six service's points are spread in the different sub-rings and they are in particular situations. Indeed in the imbricated ring solution some members of a sub-ring can be supplied by different parent routers so a part of the service will be supplied in nominal path from the same router in both solutions, but the other part will be supplied from different routers. The results of modeling are summered in Table IV.

We discover an interesting fact: the imbricated ring solution is not the optimal solution for the nominal transport. Actually this situation result from the d4 node located in the sub-ring 2. The RSTP blocking port was placed in this sub-ring considering the number of routers to join the source. We can see that d4 is able to join the source using the path d5-A3-A2-A1 and this path is shorter than the other path d3-d2-d1-c5-A1 for only one router. But this decision is an error because the availability of optical fiber is less than the microwaves links used in this study. Then it would be more logical to place the RSTP blocking port between d4 and d5 to balance in availability point of view. In this case the new engineering solution bring more stability than the actual solution, but correcting this issue on the placement of the RSTP blocking port would change this result in favor of the imbricated ring solution.

This second result is closer of the actual situation of the network, including some mistakes by deploying new infrastructures in regard to the placement of RSTP blocking ports. This result reassures again the choice to develop the new engineering solution.

TABLE IV STATE PROBABILITIES IN THE MULTIPLE SERVICE POINTS

Solution	Rings	Tree	$\Delta$
Normal	0.995651019	<b>0.99785806</b>	0.002207042
Rescue	0.004281791	<b>0.00207604</b>	0.002205751
Failure	0.000067191	<b>0.0000659</b>	0.000001291

## REFERENCES

- [1] Ghassan Semaan. Designing Networks with the Optimal Availability. OFC/NFOEC 2008. San Diego, California. USA. 2008.
- [2] Poul E. Heegaarda, Kishor S. Trivedi. Network survivability modeling. *Computer Networks*. Volume 53, Issue 8, Pages 1215–1234, 2009.
- [3] James P.G. Sterbenza, David Hutchisonb, Egemen K. Çetinkayaa, Abdul Jabbara, Justin P. Rohrer, Marcus Schöllerc, Paul Smithb. Resilience and survivability in communication networks: Strategies, principles, and survey of disciplines. *Computer Networks*. Volume 54, Issue 8, Pages 1245–1265, 2010.
- [4] Poul E. Heegaard, Kishor S. Trivedi. Survivability Quantification of Real-sized Networks including end-to-end Delay Distributions. ICSNC 2008. Sliema. Malta. 2008.
- [5] Bo Bai, Jihong Zhao, Hua Qu. A Mechanism of Maintaining the Survivability of Streaming Media Service in Overlay Networks. NAS 2013. Shaanxi. China. 2013.
- [6] Qitao Gan, Bjarne E. Helvik. Dependability Modelling and Analysis of Networks as Taking Routing and Traffic into Account. Next Generation Internet Design and Engineering. IEEE NGI 2006. Valencia. Espagne. 2006.
- [7] A. Bobbioa, L. Portinalea, M. Minichinob, E. Ciancamerlab. Improving the analysis of dependable systems by mapping fault trees into Bayesian networks. *Reliability Engineering & System Safety*. Volume 71, Issue 3, Pages 249–260, 2001.
- [8] P. Weber, G. Medina-Oliva, C. Simon, B. Lung. Overview on Bayesian networks applications for dependability, risk analysis and maintenance areas. *Engineering Applications of Artificial Intelligence*. Volume 25, Issue 4, Pages 671–682, 2012.
- [9] Abdeljabbar Ben Salem, Alexandre Muller, Philippe Weber. Dynamic Bayesian Networks in system reliability analysis. *Fault Detection, Supervision and Safety of Technical Processes*. Volume 1, Pages 444–449, 2007.
- [10] Andrew Gelman, John B. Carlin, Hal S. Stern, David B. Dunson, Aki Vehtari, Donald B. Rubin. *Bayesian Data Analysis, Third Edition*. CRC Press, Boca Raton, 2013.
- [11] D. Levi, D. Harrington. Definitions of Managed Objects for Bridges with Rapid Spanning Tree Protocol. *RFC 4318*. 2005.
- [12] B. Gleeson, A. Lin, J. Heinanen, G. Armitage, A. Malis. A Framework for IP Based Virtual Private Networks. *RFC 2764*. 2000.
- [13] M. Lasserre, V. Kompella. Virtual Private LAN Service (VPLS) Using Label Distribution Protocol (LDP) Signaling. *RFC 4762*. 2007.
- [14] B. Fenner, M. Handley, H. Holbrook, I. Kouvelas. Protocol Independent Multicast - Sparse Mode (PIM-SM) : Protocol Specification (Revised). *RFC 4601*. 2006.
- [15] B. Cain, S. Deering, I. Kouvelas, B. Fenner, A. Thyagarajan. Internet Group Management Protocol, Version 3. *RFC 3376*. 2002.